

**CS-402**

Final Term

Quiz-3,4

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**CS402**

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If the intersection of two regular languages is regular then the complement of the intersection of these two languages is \_\_\_\_\_.

**Regular**

The language of all strings partition  $\Sigma^*$  into \_\_\_\_\_ class(es).

**Two**

The language of all strings not beginning with 'b' partitions  $\Sigma^*$  into \_\_\_\_\_ distinct classes.

**Two**

If  $Q = \{xx, xyxxx\}$ , and  $R = \{xyxyxyxy, xyxyyxx\}$  then  $\text{Pref}(Q \text{ in } R) = \underline{\hspace{2cm}}$

**xyyyy**

A language ending with 'b' partitions  $\Sigma^*$  into \_\_\_\_\_ distinct classes.

**Three**

If R is regular language and Q is any language (regular/ non-regular), then Pref( \_\_\_\_\_ in \_\_\_\_\_ ) is regular.

**Q,R**

The reverse of the string sbsbb over { s, b }

**Bsbsb**

The basic approach of Myhill Nerode theorem is similar to the concept of:

**concatenation of FAs**

If there is no final state of two FAs then their \_\_\_\_\_ also have no \_\_\_\_\_ state

**union, final**

If an FA has N states then it must accept the word of length

**N**

**|**

Which of the following is a non-regular language?

**Prime**

If the FA has  $N$  states, then test the words of length less than  $N$ . If no word is accepted by this FA, then it will \_\_\_\_\_ word/words.

**accept no**

In large FA with thousands of states and millions of directed edges, without an effective procedure it is \_\_\_\_\_ to find a path from initial to final state.

**Impossible**

A problem that has decision procedure is called problem.

**Decidable problem**

Which one of the following languages is a non regular language?

**Palindrome**

Using Myhill Nerode theorem we partition  $\Sigma^*$  into distinct

does not accept any string

While finding RE corresponding to TG, we connect the new start state to the old start state by the transition labeled by

null string

According to theory of automata there are \_\_\_\_\_ types of languages

Two

While finding RE corresponding to TG, If TG has more than one final state then Introduce the new final state

The states in which there is no way to leave after entry are called

Davey John Lockers

Dead States

Waste Baskets

All of the given options

What is false about the term alphabet?

It can be an empty set.

Which of the following is used to delay the transmission of signal along the wire by one step (clock pulse)?

Delay box

### More than one

The CFG  $S \rightarrow aSb \mid ab^i$  is used to express the language

### Palindrome

A non regular language can be represented by

### None of the given options

In large FA with thousands of states and millions of directed edges, without an effective procedure it is \_\_\_\_\_ to find a path from initial to final state.

### Impossible

In polish notation, (o-o-o) is the abbreviation of \_\_\_\_\_.

### Operator - Operand - Operand

If an FA has  $N$  states then it must accept the word of length

**$N$**

Using Myhill Nerode theorem we partition  $\Sigma^*$  into distinct \_\_\_\_\_.

**classes**

The production  $S \rightarrow SS | a|b|^n$  can be expressed by Regular expression

**$(a+b)^+$**

if L1 and L2 are two regular languages, then they expressed by FAs,

**can be**

The grammatical rules which involve meaning of words are called

**Semantics**

Set of all palindromes over (a,b) is:

**Regular and finite**

The language of all strings not beginning with partitions "b" into distinct classes.

**Two**

The CFG is said to be ambiguous if there exist at least one word of its language that can be generated by production trees.

The states in which there is no way to leave after entry are called

Davey John Lockers

Dead States

Waste Baskets

All of the given options

$(a^* + b^*)^* = (a + b)^*$  this expression is \_\_\_\_\_

True

What is false about the PALINDROME LANGUAGE?

Every word is reverse of itself.

It is an infinite language.

FA can be build for it.

None of the given optio

FA is also called

DFA

Kleene star closure can be defined

Over any set of string

$\{(a + b)(a + b)\}^*$ , given RE cannot generate the string \_\_\_\_\_

bbbbbb

In an FA, when there is no path starting from initial state and ending in final state then that FA

To examine whether a certain FA accepts any words, it is required to seek the paths \_\_\_\_\_ state.

from initial to final

If  $r_1$  and  $r_2$  are regular expressions then which of the following is not regular expression

$r_1 - r_2$

Kleene star closure can be defined

Over any set of string

Which of following string(s) belongs to the language of the regular expression  $(aa^*b)^*$ ?

aabaab

According to theory of automata there are \_\_\_\_\_ types of languages

Select correct option:

Two

If  $S = \{aa, bb\}$ , then  $S^*$  will not contain

aaabbb

Every non deterministic Finite Automata can be converted into

Regular Expression

Deterministic Finite Automata

Transition Graph

All of the given options

**regular**

According to Myhill Nerode theorem, if L generates finite no. of classes then L is.....  
Select correct option:

**Regular**

Two languages are said to belong to same class if they end in the same state when they run over an FA, that state

**May be final state or not**

In  $\text{pref}(Q \text{ in } R)$  Q is ..... to (than) R.

**Not equal**

For language L defined over {a, b}, then L partitions  $\{a, b\}^*$  into ..... classes

**Distinct**

Which of the following is not a true theorem?

**Pseudo theorem**

If a regular expression contains \* then it \_\_\_\_\_ define an \_\_\_\_\_ language.

**may, infinite**

$a^n b^n$  generates the ..... language

**non regular**

Prime is a \_\_\_\_\_ language.

**non-regular**

If an effectively solvable problem has answer in YES or NO, then the solution is called \_\_\_\_\_.

**decision procedure**

To write the expression from the tree, it is required to traverse from \_\_\_\_\_.

**Left side of the tree**

If there is no final state of two FAs then their \_\_\_\_\_ also have no \_\_\_\_\_ state

**union, final**

In CFG, symbols that cannot be replaced by anything are called \_\_\_\_\_.

**terminals**

Finite Automaton (FA) must have \_\_\_\_\_ number of states while a language has \_\_\_\_\_ words.

**finite, infinite**

The language "PRIME" is an example of \_\_\_\_\_ language.

**non regular**

Using Myhill Nerode theorem we partition  $\Sigma^*$  into distinct \_\_\_\_\_.

classes

Even-Even language partitions  $\Sigma^*$  into \_\_\_\_\_ distinct classes.

four

What will be the 9's complement of the number 872?

127

In  $\text{pref}(Q \text{ in } R)$ ,  $Q$  is \_\_\_\_\_ to/than  $R$ .

Not equal

There is at least one production in CFG that has one \_\_\_\_\_ on its left side.

Non terminal

In pumping lemma theorem  $(x y^n z)$  the range of  $n$  is

$n=1, 2, 3, 4, \dots$

A language ending with 'b' partitions  $\Sigma^*$  into \_\_\_\_\_ distinct classes.

three

The operators like  $(^*, +)$  in the parse tree are considered as \_\_\_\_\_.

terminals

For a machine with N number of states, the total number of strings to be tested, defined over an alphabet of m letters, is \_\_\_\_\_.

$m^N + m^{N+1} + m^{N+2} + \dots + m^{2N-1}$

In a CFG, the non-terminals are denoted by \_\_\_\_\_.

Capital letters

If an FA accepts a word then there must exist a path from \_\_\_\_\_.

In case of Myhill Nerode theorem, if a language L partitions sigma star into distinct classes and L is also regular then L generates \_\_\_\_\_ number of classes.

finite

Which of the following is pumped to generate further strings in the definition of Pumping Lemma?

y

The complement of a regular language is also \_\_\_\_\_.

regular

An FA has same initial and \_\_\_\_\_ state, then it means that it has no \_\_\_\_\_ state.  
Select correct option:

final, initial

The product of two regular languages is \_\_\_\_\_.

If a language is regular it must generate \_\_\_\_\_ number of distinct classes.

**finite**

If L1 and L2 are regular languages then which statement is NOT true?

**L1/L2 is always regular**

If the FA has N states, then test the words of length less than N. If no word is accepted by this FA, then it will \_\_\_\_\_ word/words.

**accept no**

A problem that has decision procedure is called \_\_\_\_\_ problem.

**decidable**

If L1 and L2 are two regular languages, then they \_\_\_\_\_ expressed by FAs.

**can be**

A language that can be expressed by RE, is said to be a \_\_\_\_\_ language.

**regular**

The values of input (say a & b) do not remain same in one cycle due to

**NOT gate**

## Classes

In pumping lemma theorem ( $x y^n z$ ) the range of  $n$  is

$n=1, 2, 3, 4, \dots$

The values of input (say  $a$  &  $b$ ) do not remain same in one cycle due to

## NOT gate

The operators like ( $*$  &  $+$ ) in the parse tree are considered as

## Terminals

Even-Even language partitions  $\Sigma^*$  into \_\_\_\_\_ distinct classes.

## Four

The strings or words which do not belong to a language are called \_\_\_\_\_ of that language.

## Complement

For a machine with N number of sta., the total number of strings to be tested, defined over an alphabet of m letters, is

$$m^N + m^{N+1} + m^{N-2} + m^{2N-1}$$

Which of the following is not a true theorem?

**Pseudo theorem**

The language "PRIME" is an example of language.

**non regular**

The product of two regular languages is

**Regular**

One language can have \_\_\_\_\_ CFG(s).

**More than one**